On statistical power grid observability under communication constraints (invited paper)

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Abstract: Phasor Measurement Units (PMUs) have enabled real-time power grid monitoring and control applications realizing an integrated power grid and communication system. The communication network formed by PMUs has strict latency requirements. If PMU measurements cannot reach the control centre within the latency bound, they will be invalid for calculation and may compromise the observability of the whole power grid as well as related applications. To address this issue, this study proposes a model to account for the power grid observability under communication constraints, where effective capacity is adopted to perform a cross-layer statistical analysis in the communication system. Based on this model, three algorithms are proposed for improving power grid observability, which are an observability redundancy algorithm, an observability sensitivity algorithm and an observability probability algorithm. These three algorithms aim at enhancing the power system observability via the optimal communication resource allocation for a given grid infrastructure. Case studies show that the proposed algorithms can improve the power system performance under constrained wireless communication resources.

1 Introduction

Phasor Measurement Units (PMUs) can provide real-time power grid measurements via advanced power system and communication technologies, which improves the performance of power grid monitoring and control [1]. The PMUs are usually installed at selected buses in the power grid, which can provide measurements of both voltage and current phasor at that bus. At the same time, the communication modules associated to PMUs also form a communication network, which is synchronised by the Global Positioning Satellite (GPS). The phasor measurements are also transmitted via this communication network [2]. Since one PMU is capable to provide the information of each branch connected to that bus besides the bus itself, we can use a relatively smaller number of PMUs to monitor the whole power grid operation status. With the real-time information from PMUs deployed across the power grid, potential applications like real-time stability enhancement and state estimation [6], cyber security [7] and deployment costs [8].

From the aspect of power grid observability, PMUs show a great advantage over RTUs. It has been proposed that maintaining certain degrees of observability redundancy (OR) will be beneficial in case of PMU failures. To this end, several algorithms have been proposed, which are able to maintain the whole power system observable in the case of one or multiple PMU failures. The primary and backup (P&B) method has been proposed in [9], which consists of two independent sets of PMUs and both of them can provide full observability of the whole power grid. In [10], a local redundancy method has been proposed, which aims at guaranteeing the redundancy from the individual bus aspect. When PMU measurements are used for real-time power grid monitor, it usually requires a stringent latency performance. If the PMU measurements are not collected at the control centre within a valid latency bound, these measurements will be invalid and compromise the monitoring performance of the whole power grid. However, latency is inevitable for a practical communication system.

Compared to its wired counterpart, wireless communication technology has many advantages, such as low cost, flexibility and scalability [11]. Hence, wireless communication is playing a more and more important role in supporting the communication needs of modern grid [12]. In IEEE Standard 2030.2-2015 [13], the application of wireless technology for the communication between components within a transmission network and the operation control centre has been identified. There have been various researches addressing the wireless communication network in supporting communication between PMUs [14–17] as well as components of SCADA system [18–20]. Yet wireless communication is broadcasting in nature, which makes propagation signal prone to the influence of physical environment. The effect of channel fading will induce communication system performance fluctuation and then result in communication delay. However, the communication delay's influence on the power system observability performance as well as the inter-discipline study of the power system and communication system has not been well addressed, which is the major focus of this paper.

Communication latency is a link layer metric used in the Open Systems Interconnection (OSI) model. In practical systems, communication delay has many sources. Some latencies are fixed or bounded, such as system overheads. Others are time-varying and hard, if not impossible, to be bound. One major uncertainty contributed to this time-varying latency is due to the communication channel fading effect. However, typically latency is a metric considered in link layer but not physical layer, where the latency study is further complicated when the channel has parameters that change with time. Therefore, it requires sophisticated cross-layer analysis to study such problems. Another challenge is that, in most fading channel scenarios, it is not feasible to provide a deterministic bound for the communication delay, which is a consequence of communication performance fluctuation induced by channel fading [21]. To address these challenges, effective capacity (rate) theory is considered in this paper, which provides a cross-layer model to estimate the statistical delay bound under channel fading scenarios. The effective capacity theory is a powerful analytical tool and can be applied as a quality of service provisioning metric in various communication systems, such as cellular networks [22], multi-hop wireless networks [23] and cognitive radio networks [24]. Besides, in [25–27], the effective rates under various fading scenarios have been extensively studied.

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because of the practical situations.

In this paper, the power system observability under communication constraints is studied. The major contributions are the following:

- We propose a model to account for the power grid’s observability under constraint communication resources, which addresses the coupling effects of the power system and wireless communication system.
- The influence of wireless channel fading phenomena on the observability has been characterised using a statistical analysis and cross-layer analysis method. It provides a promising and general analysis method to bound the uncertainty effects due to wireless communication reliability.
- Three observability optimisation algorithms are proposed via an optimal communication resource allocation. The algorithms are focused on different performance metrics of observability, which are the OR, observability sensitivity (OS) and observability probability (OP).

The remaining of this paper is structured as follows. In Section 2, the PMU-based power system observability is reviewed, while the power system observability under communication constraints is studied and modelled in Section 3. The effective capacity theory is studied in Section 4, which provides a communication system cross-layer analysis framework to fill the research gap between the statistical power system observability analysis and the constraint communication resources. Then three observability improvement algorithms are proposed in Section 5, while case studies are performed in Section 6. Finally, conclusions are drawn in Section 7.

2 PMUs-based power system observability

With the PMU deployment, a lot of real-time applications have been enabled, such as state estimation, adaptive relaying and voltage instability enhancement. Compared to traditional measurement methods, PMUs are more versatile and they can provide more timely information about the power grid. Fig. 1 illustrates a typical IEEE 14-bus power system. When a bus is installed with a PMU and all branches are monitored, the PMU can measure the bus phasor voltage and outgoing current phasor information on the connected branches. The PMU can be also configured to monitor only some of the branches, but in this paper, it is assumed that all connected branches are monitored.

In an $N$-bus power system, let the binary column vector $\mathbf{x} = [x_1, x_2, \ldots , x_N]^T$ denote the PMU installation vector, where $(\cdot )^T$ is transpose operator. Its elements are given by

$$x_i = \begin{cases} 
1 & \text{if a PMU is installed at bus } i, \\
0 & \text{otherwise}. 
\end{cases}$$

Furthermore, it is assumed that the system has installed $K$ PMUs, which are labelled as $\text{PMU}_k$, where $k = 1, 2, \ldots , K$ and $K \leq N$. It can be verified that $K = \sum_{i=1}^{N} x_i$.

It is assumed that the grid topology is known as a priori for a given power grid. That is, the elements of binary network connectivity matrix $\mathbf{H}$ are known and given by

$$h_{ij} = \begin{cases} 
1 & \text{if bus } i \text{ and } j \text{ are connected or } i = j, \\
0 & \text{otherwise}. 
\end{cases}$$

A bus will be observable if at least one PMU is placed at that bus or any bus incident to it [10]. Hence, the bus observability vector $\mathbf{b}$ can be given as

$$b_i = \sum_{j=1}^{N} h_{ij} x_j. \tag{4}$$

where each element $b_i$ in the observability vector $\mathbf{b}$ indicates the number of PMUs connected to or located at the bus $i$, we have

$$b_i = \sum_{j=1}^{N} h_{ij} x_j. \tag{4}$$

It should be noticed that the observability vector $\mathbf{b}$ in (3) and its element $b_i$ in (4) are defined based on mathematical expectations. Hence the power grid is expected to be observable if $\mathbf{b} \geq 1_N$, i.e. $b_i \geq 1, \forall i$. If $b_i = 0$ for some $i$, the associated bus will not be expected to be observable. It can be seen that this grid observability model considers both power grid topology and PMU installation features. In the following section, an extended model will be proposed with the consideration of communication constraints.

3 Grid observability under communication constraints

In this section, we will further extend the observability definition in Section 2 to account for communication constraints. For a practical power system, the system statuses, such as currents, voltages and angles, would vary with time. Hence the real-time grid status monitoring of the power grid has a stringent latency requirement. To maintain real-time performance, each measurement from the PMUs will be valid within a delay bound $D_{\text{max}}$. If the measurement packages have been delayed longer than $D_{\text{max}}$ then these measurements cannot be used, which results in a compromised power grid status monitoring performance. The latency has many contributors, such as processing overheads and transmission delays, which are usually fixed values for a considered scenario. However, within wireless communication systems, the latency resulting from the channel fading effect usually varies with time and it is hard to bound. In ideal cases, the communication systems should be designed to provide a 100% guarantee that the communication delay $d_{ik}$ of PMU $i$ is smaller than $D_{\text{max}}$. However, in practice, it has been identified that it is not feasible to provide a deterministic delay bound for the communication system in most fading channel environments [21]. Hence instead, we consider the probability $0 \leq \rho_i \leq 1$ to guarantee the communication delay within a certain maximum allowed bound $D_{\text{max}}$, i.e.

$$\Pr \{d_{ik} \leq D_{\text{max}}\} \geq \rho_i. \tag{5}$$

Based on this, we can provide a statistical measure for the communication and the power system performance. It should be noted that in Section 5, we will show that providing 100% statistical guarantee is not cost effective. However, the power system performance under ideal communication scenarios can be approached via a trade-off between power system and communication system performances, which will be detailed in Section 5.

Furthermore, we define the diagonal probability matrix $\mathbf{A}_\rho = \text{diag}(\rho_1, \rho_2, \ldots , \rho_N)$, whose elements are given by

$$\mathbf{A}_\rho = \text{diag}(\rho_1, \rho_2, \ldots , \rho_N).$$

where $\mathbf{A}_\rho$ is a diagonal matrix with elements $\rho_i$.
\[ P_i = \begin{cases} p_i & \text{if PMU}_i \text{ installed at bus } i, \\ 0 & \text{otherwise.} \end{cases} \]  

For the real-time grid monitoring, if the latency of the measurements from a certain PMU exceeds \( D_{\text{max}} \), then this information will not be used. In this paper, power grid observability vector \( \mathbf{b} \) under statistical latency guarantee can be defined as follows:

\[ \mathbf{b} = \Lambda \mathbf{Q} \mathbf{X}, \]  

where \( \Lambda \) denotes the diagonal communication constraint matrix, which is defined by

\[ \Lambda = \text{diag}(Q_1, Q_2, \ldots, Q_N). \]  

where \( Q_i, i = 1, 2, \ldots, N \), is a binary random variable, which can be given by

\[ \begin{align*} 
\Pr \{Q_i = 1\} &= P_i, \\
\Pr \{Q_i = 0\} &= 1 - P_i. 
\end{align*} \]  

Therefore, the observability vector \( \mathbf{b} \) is a vector of random variables. In this paper, we focus on the observability compromised by communication performance fluctuation, where the fluctuation is due to communication channel fading effect. PMUs are installed on selected buses, which are physically and geographically separated. Hence without loss of generality, it is assumed that random variables \( Q_i \) are independent of each other. By using the fact that \( \mathbf{H} \) and \( \mathbf{X} \) are known, the expected power grid observability vector \( \mathbf{\tilde{b}} \) is given by

\[ \mathbf{\tilde{b}} = \mathbf{H} \mathbf{\Lambda}_x \mathbf{X}. \]  

The physical meaning of each element \( \mathbf{\tilde{b}}_i, i = 1, 2, \ldots, N \) of the expected grid observability \( \mathbf{\tilde{b}} \) is that the bus status information is available from an average of \( \mathbf{\tilde{b}} \) PMUs connected to the bus \( i \). If any element \( \mathbf{\tilde{b}}_i \) is smaller than 1, then it means that the observability of this bus will not be guaranteed in a statistical view, and the power grid is vulnerable to the loss of the observability of that bus.

From the power system's aspect, a full observability of the system only requires all bus observability to be one. Any extra information about that bus can be regarded as OR to that bus. The OR is not only beneficial to cope with possible PMU failures but also to improve the grid security [7]. In this paper, three different algorithms are proposed to improve the observability under a given grid infrastructure, which will be detailed in Section 5.

From (7) and (10), it can be proved that the power grid observability vector \( \mathbf{\tilde{b}} \), as well as the expected power grid observability, \( \mathbf{\tilde{b}} \) will be enhanced if the \( p_i \) for all PMUs are kept to be as close to 1 as possible. However, in practical systems, the communication system has a limited total bandwidth \( B^p \). This can be defined as a constraint for the bandwidth \( B_k \) assigned to each PMU, i.e.

\[ \sum_{k=1}^{K} B_k \leq B^p. \]  

It can be seen that the communication constraint only confines the total available bandwidth resources to each PMU, while it is the probability \( p_k \) that is directly related to the observability. Besides, the throughput of the wireless communication system is time varying due to channel fading effect. This channel fading effect on the physical layer performance will also influence the upper layers, which will result in the latencies experienced by PMUs based applications. This research gap requires a cross-layer analysis within the communication system, which will be addressed in the following section.

### 4 Cross-layer statistical delay analysis

In communication systems, Shannon channel capacity is one of the most important performance indices, which defines the maximum achievable rate for a given channel. According to Shannon channel capacity theorem, the capacity for a given channel is determined by channel bandwidth \( B \) and signal-to-noise ratio (SNR), which can be given as follows:

\[ C = B \log_2(1 + \text{SNR}). \]  

The variation of instant SNR will affect the instant system throughput in the physical layer and then results in delay at the link layer. One major source for the SNR fluctuation is channel fading, which is characterised by the physical wireless communication channel. Yet the delay aspect is not considered in the formulation above. For real-time services, such as the considered PMU based system in this paper, we require a bounded delay. If a received PMU measurement packet violates its delay bound, it will not be used and this may compromise the overall performance. It is hard or infeasible to provide a deterministic delay bound, which is due to the fact that the channel fading attenuation varies with time [21]. Hence instead, we aim to provide a statistical delay bound guarantee for the power system. In this paper, effective capacity (rate) theory is adopted, which models the cross-layer relation between the link layer behaviour and the physical channel statistical characteristics [28].

Effective capacity is the dual concept of effective bandwidth [29], and it is defined as the maximum constant rate that a fading channel can support under statistical delay constraints. The effective capacity function can be written as [27]

\[ R(\theta, B) = -\frac{1}{\theta B} \ln \mathbb{E}[e^{-\theta C}], \]  

where \( C \) denotes the instantaneous Shannon channel capacity with block transmission of duration \( T \). The parameter \( \theta \) is called QoS exponent, which is a non-negative value. The minimum required QoS exponent \( \theta \) is the value that makes the effective capacity equal to the source rate. In order to guarantee the delay performance, the QoS exponent \( \theta \) has to satisfy the constraint \( \theta \geq \theta_0 \). Moreover, when \( \theta \rightarrow 0 \), the effective capacity approaches Shannon's capacity [30].

For PMU, its effective capacity can be given as

\[ R_\text{PMU}(\theta, B_k) = -\frac{1}{\theta B_k} \mathbb{E}_\text{PMU}[e^{-\theta R(\text{PMU})}] = -\frac{1}{\theta B_k} \mathbb{E}_\text{PMU}[e^{-\theta B_k \log_2(1 + \text{SNR})}], \]  

where \( p_k \) is the average transmit SNR, which is decided by the transmit power of the communication system. The parameter \( \theta_k \) is the instantaneous channel power gain, which is determined by the fading channel characteristics.

With the definition of effective capacity and applying queuing theory, the probability of \( d_k \) within \( D_{\text{max}} \) can be given by [28]

\[ \Pr \{d_k \leq D_{\text{max}}\} = 1 - e^{-\theta_k R_\text{PMU}(\theta_k, B_k) D_{\text{max}}}. \]  

In this paper, it is assumed that PMU generates the measurements at a constant rate of \( R_k^\text{th} \). The effective capacity should be no smaller than the rate \( R_k^\text{th} \) in order to avoid unstable status, i.e.

\[ R_k(\theta_k, B_k) \geq R_k^\text{th}. \]  

By using (14)-(16), the effective capacity theory provides a cross-layer analysis framework for the study between channel fading effect, delay bound and its associated delay bound violation probability. This probability is the same one defined in (5), which affects power system observability. Hence the communication constraints’ influence on the power system observability can be characterised via the effective capacity theory. Based on this, we
can provide algorithms to improve the power system performance via the optimal communication resource allocation.

In Section 5, the effective capacity theory will be exploited as an analysis tool for improving the power grid observability. To facilitate the discussions in Section 5, we first introduce the properties of effective capacity $R_k$ here.

**Lemma 1:** The effective capacity defined in (14) has the following properties:

$$
\frac{\partial R_k(\theta_k, B_k)}{\partial B_k} \geq 0 \quad \text{and} \quad \frac{\partial R_k(\theta_k, B_k)}{\partial \theta_k} \leq 0, \forall k
$$

(17)

and $R_k(\theta_k, B_k)$ is concave in $B_k$ and $\theta_k$.

**Proof:** Taking the partial derivative of $R_k(\theta_k, B_k)$ in (14) with respect to $B_k$ and using the fact that $\theta_k$ and $B_k$ are both positive, the first term of (17) can be obtained. Then using Holder’s inequality [31, eq. (1.7.5)], it can be proved that $R_k$ is concave in $B_k$. Applying similar procedure, the partial derivative and convexity features of $R_k(\theta_k, B_k)$ can be obtained. Interested readers can refer to [21, 32] for more details. □

In theoretical communication system analysis, the Shannon capacity defined in (12) is usually used to calculate the minimum required bandwidth, which is denoted as $B_{\text{min}}$ in this paper. For a practical system, the allocated bandwidth $B_k$ has to be larger than $B_{\text{min}}$, in order to have better latency performance. If the total bandwidth is below $B_{\text{min}}$, it is for sure that the throughput of the communication system is less than the rate of the PMU measurement messages, which will lead to communication failure. Hence throughout this paper, it is assumed that $B_k > B_{\text{min}}$ has been enforced. Then with the properties of the effective capacity $R_k$, we can prove the convexity of probability $p_k$ as follows.

**Proposition 1:** The probability $p_k$ defined in (15) is convex in $B_k$ and $\theta_k$.

**Proof:** As $R_k(\theta_k, B_k)$ is concave in $B_k$ proved in Lemma 1, $-\theta_k R_k(\theta_k, B_k)$ is convex in $B_k$. Using the definition in (15), it can be proved that $p_k$ is convex in $R_k$. Following the same procedure, the convexity of $p_k$ in $\theta_k$ can be obtained. □

The probability $p_k$ is the bridge between the observability analysis (10) and the communication constraints defined in (11). Furthermore, the convexity property of $p_k$ will be useful in finding the optimal communication system configuration for the power system observability, as will be shown in the following section.

## 5 Power grid observability driven resource allocation algorithms

In power systems, the real-time measurements from PMUs are used to the monitoring of power grid status. Based on these measurements, real-time applications such as voltage stability enhancement and demand-side management can be therefore enabled. Hence it is very important to guarantee the observability of buses. In this section, three algorithms are proposed to optimise the power grid observability under communication constraints, which are aiming at different power system performance metrics, that are, the OR, OS and OP. Throughout this paper, it is assumed that the power grid and the PMU positions are known as a priori, where the focus is placed on the optimisation of the communication system to better support the services under the given configurations.

### 5.1 OR algorithm

The grid observability is critical to applications such as grid control or planning services, therefore the loss of bus status observability can result in serious consequences. The deployment of PMUs can provide real-time power grid status, which improves the power grid observability compared to traditional methods via power flow. However, the installation of PMUs will involve vast investment, which will increase the cost of the power grid operation. In fact, when a PMU is installed on a bus, it can provide information about all buses incident to this bus besides the installed bus itself [10]. By taking advantage of this feature, the PMU installation places can be selected to achieve a trade-off between cost and power grid observability [9]. With power grid topology as a priori, it is not necessary to have PMUs installed at every bus, while the desired degree of OR can be still obtained. The OR $r$ is considered in this paper, which is defined as follows [33]:

$$
r = 1 - \frac{1}{\sum B_k} \leq 1
$$

(18)

The metric $r$ gives an evaluation of the overall power network OR. For a power grid with PMU installation places as a priori, the metric $r$ is upper bounded by the ideal communication case. For a compromised communication system under resource constraints, a larger value of $r$ means that more PMUs are expected to be available to provide measurements from a statistical view. In this part, we focus on the problem of increasing OR under communication constraints, which can be formulated as follows using the effective capacity theory:

$$
\begin{align*}
\max_{B_k, \theta_k} & \quad 1 \lambda H \Lambda \theta X \\
\text{s.t.} & \quad R_k(\theta_k, B_k) \geq B_k^* \quad k = 1, 2, \ldots, K, \\
& \quad p_k = 1 - e^{-\theta_k R_k(\theta_k, B_k)} \quad k = 1, 2, \ldots, K, \\
& \quad R_k(\theta_k, B_k) = \frac{1}{B_k^*} \ln \sum_k \left( e^{-\theta_k R_k(\theta_k, B_k)} \right), \\
& \quad \sum_k B_k \leq B^*.
\end{align*}
$$

(19)

where (19) is simplified due to the fact that the power grid bus number $N$ is constant. The optimal OR is always achievable with valid $B^*$, which can be given using the following proposition.

**Proposition 2:** The power grid OR $r$ defined in (18) is convex in $B = \{B_1, \ldots, B_k\}$ and $\theta = \{\theta_1, \ldots, \theta_k\}^T$, and a feasible solution to the problem (19) always exists with every $B_k > B_{\text{min}}^*$.

**Proof:** We have the summation form of (19) as $\sum_{k=1}^N \sum_{k=1}^K \theta_k p_k$. The convexity of the redundancy $r$ and the constraints in (19) follows Lemma 1 and the convexity of $p_k$ proved in Proposition 1. The solution existence follows the fact that the domain formed by all possible $B$ is compact. □

Since the problem (19) is convex, the solution can be obtained via numerical methods such as Interior Point approach to obtain the optimal solutions. The OR algorithm has been summarised as follows.

**Algorithm 1:** OR algorithm

**I. Initialisation:**

1. obtain network connectivity matrix $H$, PMU installation vector $X$, total bandwidths $B^*$, minimum constant rate $R_k^*$, average transmit SNR $\rho_k$;
2. obtain the effective rate model $R_k(\theta_k, B_k)$ according to the fading scenario;
3. initialise bandwidth $B_k$ and QoS exponent $\theta_k$ satisfying the constraints (19);
4. initialise $n = 0$ and calculate the equivalent optimisation objective $g_n = 1\lambda H \Lambda \theta X$;

**II. Repeat:**

1. $n = n + 1$;
It can be seen that the effective capacity theory bridges not only the cross-layer analysis of the communication system but also the theoretical analysis of the power system jointly with the communication system. This cross-layer and cross-system model enables the performance optimisation of both systems, as illustrated in (19).

5.2 OS algorithm

The bus with the least expected observability within the whole grid is most vulnerable to unobservability. Hence the least bus observability can reflect the power grid’s sensitivity to losing bus observability. In this paper, we define the OS as \( b_i \), that is the least observability among all buses.

It can be seen that the buses with small observability values can be viewed as the bottlenecks to the whole power grid’s observability. From a statistical view, these buses have more influence on the whole power grid’s observability. Therefore, the power grid observability can be improved by maximising min, \( b_i \), as follows:

\[
\begin{align*}
\max_{B_0, \theta_i} & \quad \min b_i \\
\text{s.t.} & \quad R_k(\theta_i, B_0) \geq R_k^\theta, \quad k = 1, 2, \ldots, K, \\
& \quad p_k = 1 - e^{-\lambda}b_i \gamma_k \lambda_k, \quad k = 1, 2, \ldots, K, \\
& \quad R_k(\theta_i, B_0) = -1 \left( b_i \gamma_k \lambda_k \right), \\
& \quad \sum_{k=1}^{K} B_k \leq B_i.
\end{align*}
\]

Besides, it can be shown that the maximisation of the power grid observability according to the strategy above is feasible, as stated by the following proposition.

**Proposition 3:** A feasible solution to the OS algorithm defined in (20) always exists with every \( B_i > B_{i\text{min}} \).

**Proof:** Using (10), the OS of each bus \( i \) can be given by \( b_i = \sum_{k=1}^{K} b_k \gamma_k \lambda_k \). Thus the convexity of \( b_i \) follows the convexity of \( p_k \) in Proposition 1. Then the solution existence can be given by the minimax theorem [34].

The OS algorithm can follow a similar procedure as the OR algorithm, where \( g_i \) is replaced with the optimisation objective given in (20). We see that the value of min, \( b_i \), can also reflect the power grid’s reliability to the observability loss of individual buses. With a larger value of min, \( b_i \), the power grid is less sensitive to the compromised observability, which improves the power system’s reliability, at least from the viewpoint of observability.

5.3 OP algorithm

The whole power grid’s observability depends on individual bus’s observability. Hence besides considering the expected observability-based algorithms proposed in Sections 5.1 and 5.2, another algorithm is proposed in this part to provide a desired probability for the observability of individual buses above a threshold. This problem can be formulated by the optimisation of the probability that each bus’s observability is above the desired level as follows:

\[
\begin{align*}
\max_{B_0, \theta_i} & \quad \Pr \{ b \geq \lambda \} \\
\text{s.t.} & \quad R_k(\theta_i, B_0) \geq R_k^\theta, \quad k = 1, 2, \ldots, K, \\
& \quad p_k = 1 - e^{-\lambda}b_i \gamma_k \lambda_k, \quad k = 1, 2, \ldots, K, \\
& \quad R_k(\theta_i, B_0) = -1 \left( b_i \gamma_k \lambda_k \right), \\
& \quad \sum_{k=1}^{K} B_k \leq B_i.
\end{align*}
\]

The physical meaning of the desired observability level vector \( \lambda \) can be given as follows. For the case when \( \lambda = 1 \), the problem defined in (21) reduces to a statistical guarantee that every bus has unity observability. For more general cases where \( \lambda \geq 1 \), and \( \lambda \neq 1 \), the algorithm defined in (21) provides a desired statistical observability level for individual buses. It should be noted that \( \lambda \) is upper bounded by \( \lambda_{\text{max}} \) which can be calculated under an ideal communication assumption.

Here we define the solution to the problem of \( b \geq \lambda \) by the diagonal matrix \( \alpha_m \), and all the solutions form a set \( \{ \alpha_m \} \), where \( m = 1, 2, \ldots, M \). Then we can further simplify the problem defined in (21) as follows:

\[
\Pr \{ b \geq \lambda \} \equiv \Pr \{ \lambda \alpha_m X \geq \lambda \} = \sum_{m=1}^{M} \Pr \{ \alpha_m = \alpha_m \}. \tag{22}
\]

Similar to the OR algorithm discussed in Section 5.1, the optimal communication resource allocation for the maximisation of the OP is feasible, which can be given by the following proposition.

**Proposition 4:** The power grid OP defined in (21) is convex in \( B \) and \( \theta \), and a feasible solution always exists with \( \lambda \leq \lambda_{\text{max}} \) and \( B_i > B_{i\text{min}} \).

**Proof:** The desired results can be obtained following similar arguments in Proposition 2. \( \square \)

The OP algorithm can follow a similar procedure as the OR algorithm, where \( g_i \) is replaced with the optimisation objective given in (21). Note that each solution \( \alpha_m \) consists of only binary elements, namely 0 and 1. Hence when the PMU installation buses are known as a priori, the solution set \( \{ \alpha_m \} \) is readily available. Besides, if the measurements for some buses are critical information or critical to the whole power grid’s observability, we impose such buses to offer higher desired observability levels, which can be achieved by assigning corresponding elements in the parameter \( \lambda \).

It can be seen that the application of statistical analysis is a promising way to bound the uncertainties due to communication performance variations, especially when power systems and communication systems are deeply coupled with each other. Thus the methods used in the proposed algorithms are also valuable to the research of similar problems such as PMU based grid monitoring [14–17], SCADA systems [18–20] and smart meter aggregations [35, 36].

6 Case studies

In this section, the three proposed algorithms in Section 5, i.e. the OR, OS, OP algorithms, are verified using two case studies, namely IEEE 14-bus power system test case and IEEE 30-bus power system test case. These two test cases have been extensively used as standard test cases to verify power system performances [6, 10].

Here we apply the P&B method [9] for the PMU installation. The objective of P&B method is to provide the power grid with two independent PMU sets. Either P&B set is capable to provide a full observability of the whole power grid. This provides the power grid with redundancy, where the whole grid is still expected to be observable when multiple PMU’s fail within only one set. Without
loss of generality, it is assumed that all PMUs generate measurement packages at the rate of 60 kbps and the maximum allowed latency bound for these measurement packages is set to be 10 ms [37, 38]. Please note that the main topic of this paper is to improve the observability from the communication aspect, where the PMU placement is assumed as a priori knowledge which only determines the upper bound for the best observability. Moreover, it is assumed that the control centre locates at the power grid centre. The average SNR $\rho_k$ for PMU$_k$ is related to the distances between the control centre and the PMU, which are in the range of 10–15 dB in IEEE 14-bus case and 7–12 dB in IEEE 30-bus case. It is assumed that all PMU transceivers use unit transmit power. Rayleigh fading is considered as the channel fading effect for each communication channel. Under such conditions, the effective capacity under Rayleigh fading channels can be given by [27]

$$R_k(\theta_k, B_k) = \frac{1}{\theta_k T} \ln \left( \frac{\theta_k B_k T}{\ln 2 \cdot 1 - \rho_k} \right).$$

(23)

where $F \cdot \cdot$ is the generalised hypergeometric function [31].

In case of PMU communication system failure, a redundancy bandwidth is always allocated to that PMU transceiver. We adopt the bandwidth allocation algorithm similar to [39] as default algorithm, where required bandwidth for the PMUs is calculated using Shannon capacity (12) and the extra bandwidth will be evenly divided and allocated to each PMU. In our proposed algorithms, the total bandwidth is allocated according to the optimal solution of (19)–(21). For the OP algorithm, the desired statistical observability level is assumed to be $\lambda = 1_x$. Since the problems are non-linear optimisation problems, the Interior Point approach is applied to find the optimal bandwidth allocation solution.

### 6.1 IEEE 14-bus case study

In this part, the case of IEEE 14-bus power system has been considered, whose bus topology and PMU installation position have been shown in Fig. 1. The statistical probability of the communication delay associated with PMU at bus 2 has been given in Fig. 2. The Shannon capacity required for the PMU at bus 2 is 17.344 kHz under the considered scenario. It can be indicated from Fig. 2 that the latency bound will not be met with the only minimum required bandwidth. In order to counteract the latency induced communication system fluctuation, extra bandwidths are needed for the desired performance, whose quantity can be obtained via (15) and (23).

The OR under different total communication bandwidth constraints is illustrated in Fig. 3a. Here we consider the PMU loss only results from the maximum latency bound violation. Without any PMU loss, using the PMU installation position defined in Table 1, the overall grid OR can be calculated to be 0.5. It can be seen from Fig. 3a that the OR algorithm provides the best grid OR performance across different bandwidths. Specifically, with a total bandwidth of 163 kHz, the proposed OR algorithm can provide a close performance to the situation of no PMU loss, while the default algorithm requires 169 kHz to reach a similar performance.

In an ideal communication scenario, the OS algorithm provides better statistical performance over the OP algorithm when considering redundancy and sensitivity metrics. One major reason is that these two algorithms are both based on expected observability while the OP algorithm focuses on the probability performance. The OP algorithm aims at improving the probability that individual bus observability is over the desired threshold. In this case study, the desired threshold has been set to be 0.5, which equals the case that the whole power system has unity observability. It is shown in Fig. 3c that the OP algorithm provides better statistical guarantee individual bus's observability to be larger than 0.5. It can also be noticed that this performance gain is at the cost of a reduced overall OR and OS, as shown in Figs. 3a and b.

More detailed performances related to individual buses are given in Tables 2 and 3, where the total available bandwidth is 159 kHz. It is worth mentioning that, under the considered scenario, the minimum required total channel bandwidth is calculated to be 131.87 kHz using Shannon capacity theorem. However, it can be inferred from Table 2 that, with only Shannon capacity, the system observability cannot meet the requirement. Using a default algorithm, which provides each PMU with required Shannon bandwidth and evenly divides the extra bandwidth, the bus 8 is vulnerable to losing observability in the considered scenario. On the contrary, every bus observability can be statistically guaranteed by the OR, OS or OP algorithms, where the performance has been optimised for different desired performance metrics, respectively. From Table 2 as well as Figs. 3a–c, it can be seen that the proposed algorithms make better use of the extra bandwidth, to obtain performance improvements on OR, OS and OP, respectively.

### Table 1  PMU configuration

<table>
<thead>
<tr>
<th>Case</th>
<th>PMU number</th>
<th>Bus index</th>
</tr>
</thead>
<tbody>
<tr>
<td>IEEE 14-bus</td>
<td>9</td>
<td>2, 4, 5, 6, 7, 8, 9, 11, 13</td>
</tr>
<tr>
<td>IEEE 30-bus</td>
<td>21</td>
<td>1, 3, 5, 7, 8–13, 15, 17–19, 22, 24–29</td>
</tr>
</tbody>
</table>
6.2 IEEE 30-bus case study

In order to test the performance of the proposed algorithms, the IEEE 30-bus power system test case has also been investigated. The bus topology for the IEEE 30-bus power system is given in Fig. 4.

The proposed three algorithms are oriented in the optimisation of three different power system performance metrics, namely OR, OS and OP. The simulation results have been given in Figs. 5a–c. It can be seen from these figures that, the three proposed algorithms have better performance overall considered performance metrics than the default algorithm in the considered scenarios.

As illustrated in Fig. 5a, the OR algorithm provides more redundancy than the OP algorithm as well as the default algorithm. From the aspect of OR, the performance gain for the OR algorithm is slightly higher than the OS algorithm. However, this loss of the performance gain in the OS algorithm improves the power grid OS, as indicated in Fig. 5b. This is because the overall resources are constrained, which results in the situation that, the improvement of certain bus observability will be at the cost of other bus observability. Although this redundancy performance gain does not seem to be large between the OR algorithm and the OP algorithm, it should be noticed that the redundancy performance in Fig. 5a targets the whole power system performance, while the sensitivity performance in Fig. 5b targets individual buses. With constrained total resources, the improvement of the overall grid OR will be less seemingly prominent in the figures than the sensitivity performance. However, it should be noted that individual bus performances are different, as shown in Tables 2 and 3.

As can be seen in Fig. 5c, the OP algorithm improves the probability that the requirement of power system observability is met over different total communication bandwidths. It can be also seen that the performance gain is at the cost of a decrease in redundancy and sensitivity performances, as can be indicated from Figs. 5a and b.

Comparing performances between IEEE 14-bus case in Figs. 3a–c and IEEE 30-bus case in Figs. 5a–c, the three proposed algorithms provide better observability than the default algorithm when corresponding optimised metrics are considered. However, it also indicates that no single algorithm outperforms the other algorithms if all metrics are considered at the same time. The optimal algorithm depends on the considered scenario and the metric of interest.

The considered performance metrics, namely OR, OS and OP, are all formulated using a statistical approach. In theory, the best performance where an ideal communication system is considered, can be asymptotically approached. Yet from the discussion above, to improve the average performance, it has to increase the overall communication resources in an exponential way. The performance gain may be marginal even with large deployment of communication resources, especially when it is close to the best performance. Hence the results also suggest that there is a trade-off between the observability and the bandwidth. In the considered IEEE 30-bus case, the power system can reach performance similar

Table 2

<table>
<thead>
<tr>
<th>PMU bus location</th>
<th>2</th>
<th>4</th>
<th>5</th>
<th>6</th>
<th>8</th>
<th>9</th>
<th>10</th>
<th>11</th>
<th>12</th>
<th>13</th>
<th>14</th>
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</thead>
<tbody>
<tr>
<td>default</td>
<td>0.904</td>
<td>0.942</td>
<td>0.958</td>
<td>0.954</td>
<td>0.921</td>
<td>0.801</td>
<td>0.946</td>
<td>0.919</td>
<td>0.919</td>
<td>0.906</td>
<td></td>
</tr>
<tr>
<td>OR</td>
<td>0.909</td>
<td>0.942</td>
<td>0.939</td>
<td>0.921</td>
<td>0.942</td>
<td>0.894</td>
<td>0.941</td>
<td>0.894</td>
<td>0.894</td>
<td>0.894</td>
<td></td>
</tr>
<tr>
<td>OS</td>
<td>0.905</td>
<td>0.930</td>
<td>0.943</td>
<td>0.944</td>
<td>0.941</td>
<td>0.891</td>
<td>0.940</td>
<td>0.900</td>
<td>0.892</td>
<td>0.892</td>
<td></td>
</tr>
<tr>
<td>OP</td>
<td>0.904</td>
<td>0.928</td>
<td>0.943</td>
<td>0.944</td>
<td>0.941</td>
<td>0.891</td>
<td>0.940</td>
<td>0.900</td>
<td>0.892</td>
<td>0.892</td>
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</tr>
</tbody>
</table>

Table 3

<table>
<thead>
<tr>
<th>PMU bus location</th>
<th>2</th>
<th>4</th>
<th>5</th>
<th>6</th>
<th>8</th>
<th>9</th>
<th>10</th>
<th>11</th>
<th>13</th>
</tr>
</thead>
<tbody>
<tr>
<td>default</td>
<td>0.004</td>
<td>0.737</td>
<td>0.999</td>
<td>0.999</td>
<td>0.908</td>
<td>0.004</td>
<td>0.991</td>
<td>0.997</td>
<td>0.299</td>
</tr>
<tr>
<td>OR</td>
<td>0.909</td>
<td>0.942</td>
<td>0.958</td>
<td>0.954</td>
<td>0.921</td>
<td>0.801</td>
<td>0.946</td>
<td>0.919</td>
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<td>0.894</td>
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<tr>
<td>OP</td>
<td>0.904</td>
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<td>0.943</td>
<td>0.944</td>
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<td>0.891</td>
<td>0.940</td>
<td>0.900</td>
<td>0.892</td>
</tr>
</tbody>
</table>

Fig. 4 IEEE 30-bus power system with 21 PMUs

Fig. 5 Power grid observability under IEEE 30-bus power system test case

a OR performance in IEEE 30-bus case
b OS performance in IEEE 30-bus case
c OP performance in IEEE 30-bus case

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to that with ideal communication using a total bandwidth of 520 kHz.

7 Conclusion

In this paper, the power grid observability has been studied by jointly considering the power system and a wireless communication system. A corresponding analysis model has also been formulated. In order to perform the communication system cross-layer analysis, as well as to consider the channel fading effect and total bandwidth constraint, the effective capacity theory has been adopted and utilised. Based on this cross-scenario and cross-layer analysis model, three observability metrics have been formulated, namely OR, OS and OP. Then, corresponding improvement algorithms have been proposed via optimal communication resource allocation. The IEEE 14-bus and 30-bus power systems have been used in the case study to validate the performance of the three proposed algorithms. Results show that the proposed algorithms can help improve the power grid observability. Furthermore, the three proposed algorithms have the potential to be used for a trade-off between the investment needed for the communication system and the required power system performance.

8 Acknowledgment

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9 References


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